Password hacking, the de Bruijn way.

Dan Fretwell



Outline of talk





- 3 Constructing de Bruijn sequences
- 4 Magic...explained

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Magic	de Bruijn sequences	Constructing de Bruijn sequences	Magicexplained
Magic			





Here's a cool mind reading trick.



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The fool giving this talk is about to ask for five volunteers.

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Here's a cool mind reading trick.

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See...magic...

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Proper magic...

I have a deck of cards.





I have a deck of cards.

Volunteer 1, take the cards and cut the deck as many times as you want. Take the top card and pass the deck to Volunteer 2.



I have a deck of cards.

Volunteer 1, take the cards and cut the deck as many times as you want. Take the top card and pass the deck to Volunteer 2.

Volunteer $n \ge 2$, take top card from deck and pass to volunteer n + 1 until volunteer 5 has a card.

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For each of the following questions put your hand up if you satisfy the criterion:

Who has a red card?



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- Who has a red card?
- Who has a surname beginning with a letter in the first half of the alphabet? (A-M)

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- Who eats pickles on burgers?

Magic...

- Who has a red card?
- Who has a surname beginning with a letter in the first half of the alphabet? (A-M)
- Who didn't cheat on last years exams?
- Who eats pickles on burgers?
- Who is here only because they wanted to listen to Nick's talk?

How was* I able to guess all five cards?

(* replace with wasn't if necessary.)



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How was* I able to guess all five cards?

(* replace with wasn't if necessary.)

We'll find out in this talk.

Outline of talk





3 Constructing de Bruijn sequences

4 Magic...explained

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Question

How do you brute force a password of length *n* made from a finite set *X* of symbols?

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Answer

Try all $|X|^n$ possibilities!

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Answer

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If |X| or *n* is large then we really don't have time for that.

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But what if the machine lets you type continually until the correct password is entered?

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Hmmm...we would need a string of symbols that contains every $\mathbf{v} \in X^n$ as a consecutive substring at least once.

But what if the machine lets you type continually until the correct password is entered?

Hmmm...we would need a string of symbols that contains every $\mathbf{v} \in X^n$ as a consecutive substring at least once.

Obviously we could just concatenate all possibilities and enter that, but this is equivalent to the previous attack. Can we do better?

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Let's be efficient...there's no point entering a substring twice.

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Let's be efficient...there's no point entering a substring twice.

A de Bruijn sequence of order *n* for *X* is a sequence of elements of *X* such that every $\mathbf{v} \in X^n$ is a consecutive substring exactly once (allowing cycling).

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We can make some small examples easily:

If $X = \{0, 1\}$ then the following are de Bruijn sequences of order n = 1, 2, 3, 4:

01 0011 00010111 0000100110101111

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In a de Bruijn sequence each consecutive substring of order *n* has to be different and must cover all of the $|X|^n$ possibilities.

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We have proved the following:

A de Bruijn sequence of order *n* for *X*, if it exists, has length $|X|^n$.

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We have proved the following:

A de Bruijn sequence of order *n* for *X*, if it exists, has length $|X|^n$.

Note that this is much smaller than $n|X|^n$, the size of the string needed to break the password the old fashioned way!

Outline of talk





Constructing de Bruijn sequences

4 Magic...explained

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The following is non-trivial!

Theorem

If *X* is a finite set and $n \ge 1$ then there exists a de Bruijn sequence of order *n* for *X*.

In fact, if k = |X| then there are $\frac{(k!)^{k^{(n-1)}}}{k^n}$ of them.

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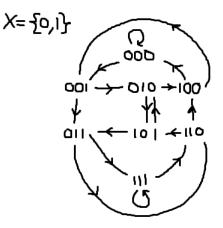
In fact, if
$$k = |X|$$
 then there are $\frac{(k!)^{k^{(n-1)}}}{k^n}$ of them.

Question

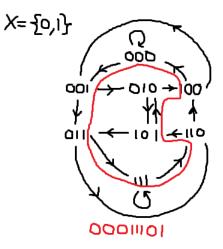
How do we make one?

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We can make de Bruijn sequences by finding Hamiltonian paths in certain graphs.



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In the case where $X = \mathbb{Z}/p\mathbb{Z}$ we can construct de Bruijn sequences recursively.

Let $f(t) = t^n + A_{n-1}t^{n-1} + ... + A_1t + A_0 \in (\mathbb{Z}/p\mathbb{Z})[t]$ be irreducible. Then the recursion:

$$x_{m+n} \equiv A_{n-1}x_{m+n-1} + \dots + A_1x_{m+1} + A_0x_m \mod p$$

gives a de Bruijn sequence of order *n* for any non-zero initial vector $\mathbf{v} = (x_0, x_1, ..., x_{n-1}) \in (\mathbb{Z}/p\mathbb{Z})^n$.

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Example

If p = 2 then $t^3 + t + 1$ is an irreducible polynomial of degree 3 over $\mathbb{Z}/2\mathbb{Z}$.

Starting with seed 001 we get the sequence:

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If p = 2 then $t^3 + t + 1$ is an irreducible polynomial of degree 3 over $\mathbb{Z}/2\mathbb{Z}$.

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001<mark>0111</mark>

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Outline of talk





3 Constructing de Bruijn sequences



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So how did the magic trick work? It should be clear that only the red card question was relevant.

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But surely that isn't enough information...there are only $2^5 = 32$ possible answers to that question but many more possible 5-tuples of cards.

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Confession 1

The deck contained only 32 cards, the ones with numerical value up to 8.

So how did the magic trick work? It should be clear that only the red card question was relevant.

But surely that isn't enough information...there are only $2^5 = 32$ possible answers to that question but many more possible 5-tuples of cards.

Confession 1

The deck contained only 32 cards, the ones with numerical value up to 8.

Confession 2

The deck was rigged!

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We can encode the cards in our deck as binary strings of length 5, two bits for the suit and three for the number:

For example:

 $5H \rightarrow 11101 \qquad \text{AC} \qquad 00001 \qquad 8S \rightarrow 01000$

00001001011001111100011011101010

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Here is a de Bruijn sequence of order 5 for $X = \{0, 1\}$:

00001001011001111100011011101010

I can use this to set up a deck:

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AC

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I can use this to set up a deck:

AC, 2C

00001001011001111100011011101010

I can use this to set up a deck:

AC, 2C, 4C



00001001011001111100011011101010

I can use this to set up a deck:

AC, 2C, 4C, AS



00001001011001111100011011101010

I can use this to set up a deck:

 $AC, 2C, 4C, AS, ..., \textcolor{red}{8C}$



00001001011001111100011011101010

I can use this to set up a deck:

 $AC, 2C, 4C, AS, ..., \textcolor{red}{8C}$

Cutting the deck corresponds to cycling the sequence...which doesn't change the de Bruijn property.

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00001001011001111100011011101010

I can use this to set up a deck:

 $\mathsf{AC}, \mathsf{2C}, \mathsf{4C}, \mathsf{AS}, ..., \textcolor{red}{\mathsf{8C}}$

Cutting the deck corresponds to cycling the sequence...which doesn't change the de Bruijn property.

Note that the 1's in the sequence translate into red cards. So knowing who has a red card gives a binary sequence of length 5 corresponding to what the first card drawn was!

How did I know the other four cards?



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How did I know the other four cards?

The de Bruijn sequence (x_n) in the previous slide is generated by the recursion $x_{n+5} \equiv x_n + x_{n+2} \mod 2$ (with $\mathbf{v} = 00001$).

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So once I know what the first card drawn was, I can generate the binary sequences corresponding to the next four cards! For example:

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For example:

 $\textbf{010110} \rightarrow \textbf{010110}$

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For example:

 $01011001 \to 01011001$

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For example:

 $010110011 \rightarrow 010110011$

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For example:

 $\textbf{010110011} \rightarrow \textbf{3S}$

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For example:

 $\textbf{010110011} \rightarrow \textbf{3S}, \textbf{6D}$

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For example:

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So once I know what the first card drawn was, I can generate the binary sequences corresponding to the next four cards!

For example:

 $\texttt{010110011} \rightarrow \texttt{3S}, \texttt{6D}, \texttt{4S}, \textbf{AH}$

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For example:

 $\texttt{010110011} \rightarrow \texttt{3S}, \texttt{6D}, \texttt{4S}, \texttt{AH}, \texttt{3D}$

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Thanks for listening